

MAHARASHTRA STATE BOARD OF TECHNICAL EDUCATION (Autonomous)

(ISO/IEC - 27001 - 2005 Certified)

MODEL ANSWER

SUMMER-17 EXAMINATION

Subject Title: System Programming

Subject Code:

17634

Important Instructions to examiners:

- 1) The answers should be examined by key words and not as word-to-word as given in the model answer scheme.
- 2) The model answer and the answer written by candidate may vary but the examiner may try to assess the understanding level of the candidate.
- 3) The language errors such as grammatical, spelling errors should not be given more Importance (Not applicable for subject English and Communication Skills.
- 4) While assessing figures, examiner may give credit for principal components indicated in the figure. The figures drawn by candidate and model answer may vary. The examiner may give credit for any equivalent figure drawn.
- 5) Credits may be given step wise for numerical problems. In some cases, the assumed constant values may vary and there may be some difference in the candidate's answers and model answer.
- 6) In case of some questions credit may be given by judgement on part of examiner of relevant answer based on candidate's understanding.
- 7) For programming language papers, credit may be given to any other program based on equivalent concept.

Q.	Sub	Answer	Marking
No.	Q. N.		Scheme
1.		Solve any FIVE:	20Marks
	(1)	List the four components of system programming.	4M
	Ans:	Assembler: The program known as assembler is written to automate the translation of	(Each
		assembly language to machine language. Input to the language is called as source	Component:
		program and output of assembler is machine language translation called as object	1 mark, any
		program.	four
		$ALP \rightarrow ASSEMBLER \rightarrow Machine Language equivalent + Information required$	Components)
		by the loader	
		Loader: Loader is a system program which places program into the memory and prepares for execution. Loading a program involves reading the contents of the	
		executable file containing the program instructions into memory, and then carrying out other required preparatory tasks to prepare the executable for running. Once loading is complete, the operating system starts the program by passing control to the loaded program code. E.g. Boot Strap loader.	
		Macro: A macro is a rule or pattern that specifies how a certain input sequence (often a sequence of characters) should be mapped to a replacement output sequence (also often a sequence of characters) according to a defined procedure. The mappings process that instantiates (transforms) a macro use into a specific sequence is known as macro expansion. A facility for writing macros may be provided as part of a software application or as a part of a programming language. In the former case, macros are used to make tasks using the application less repetitive. In the latter case, they are a tool that allows a programmer to enable code reuse or even to design domain-specific languages	



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	MACR MEND Compi source comput code).T prograr Linker object n a progr					
(2)	What a	re the g	goals of system softwa	ire?		4M
Ans:	Goal of 1) 2) 3) 4) 5) 6)	(Each Goal: 1 mark, any four goals)				
 (3)	Compa	re bina	ry search and linear s	search.		4M
Ans:		SR. NO 1 2 3 4 5 6 7 7	PARAMETERS Time Complexity Best case time Prerequisite for an array Can be implemented on Algorithm type Usefulness No of Comparison	BINARY SEARCH O (log2 N) O(1) Center Element Array must be in sorted order Cannot be directly implemented on linked list Divide and conquer in nature tricky algorithm tricky algorithm of comparisons are required	LINEAR SEARCHO(N)O(1)FirstElementNo prerequisiteArray Linked listIterative natureEasy to useNumber comparisons are less	(Each Comparison Point: 1 mark, any four Comparison points)



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<u>Subj</u> ect	Title: System Programming Subject Code: 1763	4
(4)	Define the terms:	4M
	i. Allocation ii. Relocation iii. Linking iv. Loading	
Ans:	i. Allocation: Allocate the space in the memory where the object programs can be loaded for execution. It allocates the space for program in the memory, by calculating the size of the program. This activity is called allocation.	(Each term Definition: 1 mark)
	ii. Relocation: Adjust the address sensitive instructions to the allocated space. There are some address dependent locations in the program, such address constants must be adjusted according to allocated space, such activity done by loader is called relocation.	
	iii. Linking: Resolving external symbol reference. It resolves the symbolic references (code/data) between the object modules by assigning all the user subroutine and library subroutine addresses. This activity is called linking.	
	iv. Loading: Placing the object program in the memory in to the allocated space. Finally it places all the machine instructions and data of corresponding programs and subroutines into the memory. Thus program now becomes ready for execution, this activity is called loading.	
(5)	Explain four basic operations of macroprocessor.	4M
Ans:	 The 4 basic task of Macro processor is as follows:- 1) Recognize the macro definitions. 2) Save the Macro definition. 3) Recognize the Macro calls. 4) Perform Macro Expansion. 1) Recognize the Macro definitions: - A microprocessor must recognize macro definitions Identified by the MACRO and MEND pseudo-ops. When MACROS and MENDS are nested, the macro processor must recognize the nesting and correctly match the last or outer MEND with the first MACRO. 2) Save the Macro definition: - The processor must store the macro instruction definitions which it will need for expanding macro calls. 3) Recognize the Macro calls: - The processor must recognize macro call that appear as operation mnemonics. This suggests that macro names be handled as a type of opcode. 4) Perform Macro Expansion: - The processor must substitute for macro 	(Description of each Operation: 1 mark)
	definition arguments the corresponding arguments from a macro call, the resulting	
(6)	symbolic text is then substituted for the macro call.	AM
		TIVI
Ans:	The subroutines of a program are needed at different times. For e.g. Pass 1 and pass 2 of an assembler are call other subroutine it is possible to produce an overlay structure that identifiers mutually exclusive subroutines. In order for the overlay structure to work it is necessary for the module loader to the various procedures as they are	(Description: 3 marks, Diagram: 1



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2.		Solve any FOUR:	16Marks
	(1)	Describe the concept of compile & go loader.	4M
	Ans:	One method of performing the loader functions is to have the assembler run in one part of memory and place the assembled machine instructions and data, as they are assembled, directly into their assigned memory locations. As a usual practice one method of performing the loader functions is to have to assemble run in one part of memory and place the assembled machine instructions and data they are assembled, directly into their assigned memory locations. When the assembly is completed the assembler causes transfer to the instruction of the program. This is a simple solution, involving no extra procedures. It is used by the WATFOR FORTRAN compiler and several other language processors. Such a loading scheme is commonly called "compile-and-go" or "assembler – and –go". It is relatively easy to implement. The assembler simply places the code into core, and the "loader" consists of one instruction that transfers to the starting instruction of the newly assembled program.	(Description:3 marks, Diagram: 1 mark)
		Disadvantages:	
		1. A portion of memory is wasted because the core occupied by the assembler is unavailable to the object program.	
		2. It is necessary to retranslate (assemble) the user's program code every time it is run.	
		3. It is very different to handle multiple subroutines in assembly language and another subroutine in any programming language.	
		This last disadvantage makes it very difficult to produce orderly modular programs in the design of assemblers. For example assembler is one of the type of compile and go loader which can be depicted in the following figure:	
		Source pgm Compile - and - Go translator Control Assembler {Compile - and Go Assemble - and Go	



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(2)	Draw the foundation of system software.	4M
Ans:	People Application programming Compilers Assemblers Macroprocessors Loaders Text editors Debugging aids Searching and sorting I/O programs File systems Scheduler Libraries Memory management	(Correct Diagram: 4 marks)
(3)	Explain bottom-up parsing technique.	4M
Ans:	Bottom-up parsing starts from the leaf nodes of a tree and works in upward direction till it reaches the root node. Here, we start from a sentence and then apply production rules in reverse manner in order to reach the start symbol. The image given below depicts the bottom-up parsers available.	(Description: 2 Marks, Types: 2 Marks)
	 Shift-Reduce Parsing: Shift-reduce parsing uses two unique steps for bottom-up parsing. These steps are known as shift-step and reduce-step. Shift step: The shift step refers to the advancement of the input pointer to the next input symbol, which is called the shifted symbol. This symbol is pushed onto the stack. The shifted symbol is treated as a single node of the parse tree. Reduce step: When the parser finds a complete grammar rule RHS and replaces it to LHS, it is known as reduce-step. This occurs when the top of the stack contains a handle. To reduce, a POP function is performed on the stack which pops off the handle and replaces it with LHS non-terminal symbol. 	



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(4)	 LK Parser: The LK wide class of context technique. LR parser scanning of the input reverse, and k denote There are three widel SLR1 – Simple LR P Works on sma Few number o Simple and fas LR1 – LR Parser: Works on com Generates larg Slow construct LALR1 – Look-Ahe Works on inter Number of states are 	parser is a not t free gramm is are also kn is stream; R s is the number y used algorit arser: llest class of f states, hence it construction plete set of I e table and lat tion ead LR Parse rmediate size e same as in S with example	on-recurs har which nown as tands for of look ithms ava- gramma ce very su n LR1 Gran arge num r: of gram SLR1 e.	sive, shift-reduce, h makes it the mo LR parsers, wher r the construction ahead symbols to ailable for constru r mall table mmar aber of states	bottom-up parser. It uses a ost efficient syntax analysis re L stands for left-to-right of right-most derivation in make decisions. cting an LR parser:	4M
 (4)						41v1
	is then assigned to a have been distributed repeated until no mo There are serious disa 1) It takes two separa 2) It requires a lot of The average time req is the maximum key is N *P.	bucket unique d the buckets re digits are advantages to te processes, extra Storage uired for the size & P is t	uely dep s items a left. A n o using it a separa e for the s sort is (the radix	end on the value of are merged in ord number system of t internally on a di- ation and a merge buckets. (* N *log P(K)) v of the radix sort.	of the digit. After all items ler and then the process is base P requires P buckets. gital compiler where N is the table size, K The extra storage required	2 marks, Example:2 marks)
	Original table	First Distribution	Merge	First Distribution	Final merge	
	13	0)	31	0) 01, 02, 05,	02	
	05	1) 01, 31, 11, 21	11	1) 11, 13, 16, 19	05	
	27	2) 02	21	2) 21, 26, 27	09	
	01	5) 15 4)	02	5) 51 4)	11	
	20	+) 5) 05	05	4) 5)	16	
	16	6) 26, 16	26	6)	19	
	02	7) 27	16	7)	21	
	09	8)	27	8)	26	
	11	9) 19, 09	19	9)	27	
	21	09		31		
	T Separate, based	on last digit	1 S	Separate, based on first	digit	



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(5)	Explain the concept of subroutine linkages in loader.	4 M
	To understand the concept of subroutine B is made. The subroutine B is not written in the program "A" call to subroutine B is made. The subroutine B is not written in the program segment of A, rather B is defined in some another program segment C" Nothing is wrong in it. But from assembler's point of view while generating the code for B, as B is not defined in the segment A, the assembler cannot find the value of this symbolic reference and hence it will declare it as an error. To overcome problem, there should be some mechanism by which the assembler should be explicitly informed that segment B is really defined in some other segment C. Therefore whenever segment B is used in segment A and if at all B is defined in C, then B must -be declared as an external routine in A. To declare such subroutine as external, we can use the assembler directive EXT. Thus the statement such as EXT B should be added at the beginning of the segment A. This actually helps to inform assembler that B is defined somewhere else. Similarly, if one subroutine or a variable is defined in the current segment and can be referred by other segments. This overall process of establishing the relations between the subroutines can be conceptually called a_subroutine linkage. For example MAIN START EXT B	(Description: 4 marks, any relevant description shall be considered)



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format	
Columns	Contents
1	Hexadecimal byte X'02'
2-4	Characters TXT
5	Blank
6-8	Relative address of first data byte
9-10	Blanks
11-12	Byte Count (BC) = number of bytes of information in cc. 17-72
13-16	Blank
17-72	From 1 to 56 data bytes
73-80	Card sequence number
-	1

RLD card

The RLD cards contain the following information 1. The location and length of each address constant that needs to be changed for relocation or linking. 2. The external symbol by which the address constant should be modified. 3. The operation to be performed.

RLD card format:

Columns	Contents
1	Hexadecimal byte X'02'
2-4	Characters RLD
5-18	Blank
19-20	Relative address of first data byte
21	Blanks
22-24	Byte Count (BC) = number of bytes of information in cc. 17-72
25-72	Blank
73-80	Card sequence number

END card

The END card specifies the end of the object deck.

END card format:



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1 Hexadecimal byte X'02' 2.4 Characters END 5 Blank 6-8 Start of execution entry (ADDR), if other than beginning of program 9-72 Blanks 73-80 Card sequence number (3) State four functions of compiler. Ans: 1. Recognizing basic elements like variable, keyword, operators. 2. Recognizing combination of elements as a syntactic unit and interpret their meaning Storage allocation for variables 4. Generate the equivalent object code. (for the terms: (i)Binder (ii). Dynamic loader (iii). Linking editor (iv). Overlays. (Ea Ans: 1. Binders: A binder is a program that performs the same functions as the direct-linking loader in "binding" subroutines together, but rather than placing the relocated and linked text directly into memory, it outputs the text as a file or card deck. This output file is in a format ready to be loaded and is typically called aload module. 2) Dynamic Loader: The portion of the loader that actually intercepts the "calls" and loads the necessary procedure is called the overlay supervisor or simply the flipper. This overall scheme is called dynamic loading – on call (LOCAL). 3) Linkage Editor: A more sophisticated binder, called a linkage editor, can keep track of then relocated and thereby loaded anywhere in core. 4) Overlay: An overlay is a part of a program (or software package) which has the same load origin as some other p			Columns Contents	
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 4) Overlay: An overlay is a part of a program (or software package) which has the same load origin as some other part(s) of the program. Overlays are used to reduce the main memory requirement of a program. (5) Write algorithm for syntax analysis phase of compiler. 		ceep ıble,	3) Linkage Editor: A more sophisticated binder, called a linkage editor, can track of then relocation information so that the resulting load module, as an ensen can be further relocated and thereby loaded anywhere in core.	
(5) Write algorithm for syntax analysis phase of compiler.		the the	4) Overlay: An overlay is a part of a program (or software package) which has same load origin as some other part(s) of the program. Overlays are used to reduce main memory requirement of a program.	
	4M		Write algorithm for syntax analysis phase of compiler.	(5)
Ans: Syntax analysis Phase: Algorithm (Fe 1. Reduction are checked repeatedly for match between the field which old top of stack and the actual top of stack , until match is found (Fe 2. If match is found, action routines specified in the action field are executed. (Fe 3. After syntax analyzer get control back to it, it modifies the top of stack with new top	ur os: 4 ·ks)	tack	 Syntax analysis Phase: Algorithm 1. Reduction are checked repeatedly for match between the field which old top of s and the actual top of stack , until match is found 2. If match is found, action routines specified in the action field are executed. 3. After syntax analyzer get control back to it, it modifies the top of stack with new 	Ans:



SUMMER-17 EXAMINATION

ubject	Title: System Prog	ramming			Subject Code:	17634	4
	of stack 4. Step 1 is repeated	d starting with	the reduc	tion specified in t	he next reduction	field.	
(6)	Explain the struct (ALA) with examp	ure of Macro ple.	Definitio	n Table (MDT) a	and Argument Li	st Array	4M
Ans:	MDT: - MDT is us parameter is replace definition. In pass 2 Structure of MDT	sed to save the ed by the inde 2 MDT is used	e macro de ex notation l for perfor	efinition along wi n. In pass 1 MDT rming macro expa	ith MEND statem T is used to save t ansion	ent. Each he macro	(MDT wi Example: marks, ALA with
		INDEX		MACRO DEF BYT	INITION TABLE ES/ENTRY	2 80	Example: marks)
	I AL A'-IT IS USED TOT	- simnlitvinσ fl	ne narame	ier renigeement n	rocedure in pass	I ALA 1S	
	used to replace the is used to replace the Dummy arg in the MDT Structure of ALA	formal parame ne index notati macro definitions index marker	eter by thr ions by thr on are plac	ree respective inde ree actual paramet ced with positiona	ex notations In pass ters. al indicators when d by #i	ss 1 ALA stored in	
	<pre>weight is used to replace the is used to replace the Dummy arg in the MDT Structure of ALA # symbol is used as INDEX</pre>	formal parame ne index notati macro definitions index marker	eter by thr ions by thr on are place	arg represented ARGUMENT (8 ENTRY)	al indicators when to by #i	ss 1 ALA stored in	
	<pre>wide is used to replace the is used to replace the is used to replace the MDT Structure of ALA # symbol is used as INDEX</pre>	formal parame ne index notati macro definitions index marker	eter by thr ions by thr on are place ; ,ith dumm	arg represented ARGUMENT (8 ENTRY)	al indicators when by #i	ss 1 ALA stored in	
	ALA:-It is used for used to replace the is MDT Structure of ALA # symbol is used as INDEX Example:-	formal parame ne index notati macro definitions index marker	eter by thr ions by thr on are place • ,ith dumm	arg represented ARGUMENT (8 ENTRY)	al indicators when	ss 1 ALA stored in	
	<pre>is used to replace the is used to replace the is used to replace the MDT Structure of ALA # symbol is used as INDEX Example:-</pre>	formal parame ne index notati macro definiti i s index marker	eter by thr ions by thr on are place , ith dumm	arg represented ARGUMENT (8 ENTRY)	al indicators when BYTES PER	ss 1 ALA stored in	
	ALA:-It is used for used to replace the is used to replace the Dummy arg in the MDT Structure of ALA # symbol is used as INDEX Example:-	formal parame ne index notati macro definitions index marker Solutions Solutions Mage Solutions	eter by thr ions by thr on are place • ,ith dumm • ,ith dumm ource code facro LAB INCR	arg represented ARGUMENT (8 ENTRY)	al indicators when	ss 1 ALA stored in	
	<pre>is used to replace the is used to replace the is used to replace the MDT Structure of ALA # symbol is used as INDEX Example:-</pre>	formal parame ne index notati macro definitie : s index marker So M & &	eter by thr ions by thr on are place to are place to are place to are place to are place the placet the	ee respective inde ee actual paramet ced with positiona ny arg represented ARGUMENT (8 ENTRY) &A1,&A2,&A3	al indicators when BYTES PER	ss 1 ALA stored in	
	ALA:-It is used for eplace the is used to replace the is used to replace the MDT Structure of ALA # symbol is used as INDEX Example:-	formal parame ne index notati macro definiti s index marker Solution Solution Marker A	 ine parameter by three ions by three ions by three ions by three ions is the ion are place. , ith dumm , ith dumm	ARGUMENT (8 ENTRY)	al indicators when BYTES PER	ss 1 ALA stored in	
	ALA:-It is used for eplace the is used to replace the is used to replace the MDT Structure of ALA # symbol is used as INDEX Example:-	formal parame ne index notati macro definitions s index marker So Ma &l &l &l &l &l &l &l &l &l &l &l &l &l	eter by thr ions by thr on are place to are placeto to	<pre>% arg represented ARGUMENT (8 ENTRY)</pre>	al indicators when BYTES PER	ss 1 ALA stored in	



MODEL ANSWER

SUMMER-17 EXAMINATION

Subject Title: System Programming

Subject Code:

			MNT	[
			INDE	EX N	AME	MDT POINTER(4 BYES)		
					VCRbbbb			
						131		
			Pass 2 (ALA) LOOP INCR D	DATA1,DA1	ГА2,DATA3	~ • •		
			INDEX			ARGUMENT(8 BYTES PER I	ENTRY)	
			0			"LOOPbbbb"		
			1			"DATA1bbb"		
			2		`	"DATA2bbb"		
			3			"DATA3bbb"		
		b denotes :- blank						
		MDT.						
			IND	ЭEX	INDEX DEFINI	MACRO TION TABLE 80		
					BYTES/	ENTRY	_	
		15 &LAB INCR &A1,&A2,&A3						
			1	6	#0 A 1,#	⁴ 1		
			1	7	A 1,#2			
			1	8	A 1,#3			
			1	9	MEND			
4.		Solve any FOUR:						16Marks
	(1)	Explain working o	4M					
	Ans:	Loader avoids post changed and perfor loader allows man assembles each pro	sible reas m the tas y proced cedure so	ssembl sks of a ure seg egment	ing of al allocatior gments y t indepen	Il subroutines when and linking for the vet only one data su idently and passes of	a single subroutine is programmer. The BSS egment. The assembler n to the loader the text	(Working: 4 marks)
		and information as	to reloca	ition ai	nd inter s	segment reterences.	The o/p of a relocating	



SUMMER-17 EXAMINATION

bject Title ass pro tra by inf veo wo tra Th ins Th ins Th the all 2) Dr .ns: 1.	le: System Progra ssembler using a E rogram it reference ansfer vector that c y the source progra formation such as ector position. Aft ould load each su ansfer instruction t he BSS loader sch istruction format. he relocation bit so problem of link location. raw format of MC	Amming SSS schen es. For eac consist of am. The a sthe leng ter loading ubroutine to the corr neme is of olves the p king and DT, POT,	ne is the ch source addresses assemble g the tex identifie respondin ther used problem of the prog	object p program s contain r would e entire p at and th d in the g subrou on com of reloca ram leng	rogram n the sing na also p programe trans trans time in piler tion, t gth in	Sul m and inform assembler of ames of the s provide the l um and the l nsfer vector. fer vector. n each entry with a fixed the transfer v formation so	bject Code: nation about p a text pref bubroutines r oader with a length of th into core, t It would the in the transf length direct vector is used olves the pr	17634 t all other fixed by a eferenced additional e transfer he loader e place a fer vector. ct address d to solve roblem of	1
ass pro tra by inf veo wo tra Th ins Th the allo 2) Dr ns: 1.	ssembler using a E rogram it reference ansfer vector that c y the source progra formation such as ector position. Aft rould load each su ansfer instruction t he BSS loader sch istruction format. he relocation bit sc ie problem of link location. raw format of MC	BSS schen es. For eac consist of am. The a s the leng ter loading abroutine to the corr neme is of olves the p cing and DT, POT,	ne is the ch source addresses assemble: th of the g the tex identifier respondin ther used problem of the prog	object program e program s contain r would e entire p at and th d in the g subrou of reloca ram leng	rogram n the a ing na also p program trans time in piler tion, t gth in	m and inform assembler o/ ames of the s provide the l um and the l nsfer vector offer vector. n each entry with a fixed the transfer v formation so	nation about p a text pref ubroutines r oader with a length of th into core, t It would the in the transf length direct vector is used olves the pr	t all other fixed by a eferenced additional e transfer he loader e place a fer vector. ct address d to solve roblem of	
.ns: 1.			, 51, D1		e or p	ass-1 of asse	mprer.		4 M
	Machine Op-code Mnemonic opcode 4 Bytes	e Table (I Binary (1 Byte	MOT):- Opcode	Length Bit	2	Instruction Format 3 Bi	Reserve	ed 3Bit	(Each Databas format: mark)
Ps	Pseudo-Opcode Table (POT):- Pseudo-Opcode 5 Bytes Address of Function 3 Bytes =24 bit address						Bytes		
Sy	ymbol Table(ST) :	:-	Value 4 D	Bytes	Leng	gth 1 Bytes	Relocation	1 Bytes	



MODEL ANSWER

SUMMER-17 EXAMINATION

Subject Title: System Programming

Subject Code:

	Availability Designated relative-address Indicator Contents of base register (1-byts) (3-bytes = 24-bit address) (character) (hexadiccimal) 1 "N" 2 "N" 14 "N" 15 "Y" 16 "Y" 17 00 00 00 18 entries 19 "Y" 00 00 00 10 "Y" 00 00 14 "N" "Y" 00 00 00 00 00 15 "Y" Code+ Availability Y Y register specified in USING pseudo-op N N register never specified in USING pseudo-op subsquartity made unavailable by the DROP pseudo-op Pseudo-op subsquartity made unavailable by the DROP pseudo-op Pseudo-op pseudo-op	
(3)	Draw macro instruction structure.	4M
Ans:	MACRO ← Start of definition Macro name Body of macro	(Structure: 4 marks)
	(3) Ans:	(3) Draw macro instruction structure. Ans: MACRO ← Start of definition → Body of macro



MODEL ANSWER

SUMMER-17 EXAMINATION

<u>Subj</u> ect	Title: System Programming Subject Code: 1	7634	
	MEND End of macro		
	E.g:-		
	MACRO		
	INCR		
	AI, DATA		
	A2, DATA		
	MEND		
	INCR		
	•		
	END		
(4)	Explain the structure of:		4M
	i. Identifier Table ii. Matrix Database of compiler.		
Ans:	IDENTIFIER TABLE: Created by lexical analysis to describe all identifiers us the source program. There is one entry for each identifier. Lexical analysis created entry and places the name of the identifier into that entry. Since in many lange identifiers may be from 1 to 31 symbols long. The lexical phase may enter a point the identifier table for efficiency of storage. The pointer points to the name in a tal names. Later phases will fill in the data attributes and address of each identifier.	ed in (I es the T uages m ter in D ble of ta m	dentifier able : 2 arks, Matrix atabase ble: 2 arks)
	Name Data Attributes Address		
	MATRIX: The compiler uses the matrix, which is the intermediate form of a representation of the past tree. Each matrix entry has one operator and two operand intermediate result store in temporary variable.	linear rands	



Subject Title: System Programming

MAHARASHTRA STATE BOARD OF TECHNICAL EDUCATION (Autonomous) (ISO/IEC - 27001 - 2005 Certified)

MODEL ANSWER

SUMMER-17 EXAMINATION

Subject Code:

MATRIX OPERATOR OPERAND 1 OPERAND 2 RESULT LINE NO Image: An image							
(5) Explain Random Entry Scarching. 4M Ans: All Binary Scarch algorithms, which arc fast, but can only operate on tables that arc ordered and packed, i.e. tables that will have adjacent items ordered by keywords. Such search procedures may therefore have to be used in conjunction with a sort algorithm which both orders and packs the data. Actually, it is unnecessary for the table to be ordered and packed to achieve good speed in searching. This is also possible to do considerably better with an unpacked, unordered table, provided it sparse, i.e. the number of storage spaces allocated to it exceeds the number of items to be stored. It is observed that the address calculation sort gives good results with a sparse table. However, having to put elements in order slows down the process. A considerable improvement can be achieved by inserting element in a random (or pseudo-random) way. The random entry number K is generated from the key by methods similar to those used in address calculation. If the K the some other cell must be found for the insertion. The first problem is the generation of a random number from the key. It is to design a procedure that will generate pseudo-random, consistent table positions for keywords by the table length N and use the remainder. This scheme works well as long as N and the key size (32 bits in case) have no common factors. For a given group of M keywords the remainders should be fairly evenly distributed over) (N-1). (6) Explain Machine dependent optimization phase of compiler with example. 4M Ans: Machine dependent optimization phase of compiler with example. 4M Ans: Machine dependent optimization phase of compiler with example. 4M Ans:		MATRIX LINE NO	OPERATOR	OPERAND 1	OPERAND 2	RESULT	
(5) Explain Random Entry Searching. 4M Ans: All Binary Search algorithms, which are fast, but can only operate on tables that are ordered and packed, i.e. tables that will have adjacent items ordered by keywords. Such search procedures may therefore have to be used in conjunction with a sort algorithm which both orders and packs the data. Actually, it is unnecessary for the table to be ordered and packed to achieve good speed in searching. This is also possible to do considerably better with an unpacked, unordered table, provided it is sparse, i.e. the number of storage spaces allocated to it exceeds the number of items to be stored. It is observed that the address calculation sort gives good results with a sparse table. However, having to put elements in order slows down the process. A considerable improvement can be achieved by inserting element in a random (or pseudo-random) way. The random entry number K is generated from the key by methods similar to those used in address calculation. If the K the some other cell must be found for the instrotion. The first problem is the generation of a random number from the key. It is to design a procedure that will generate pseudo-random, consistent table positions for keywords. One fairly good prospect for four character EBCDIC keywords is to simply divide the keyword by the table length N and use the remainder. This scheme works well as long as N and the key size (32 bits in case) have no common factors. For a given group of M keywords the remainders should be fairly evenly distributed over) (N-1). (M Ans: Machine dependent optimization phase of compiler with example. 4M Ans: Machine dependent optimization phase of compiler with example. 4M Ans: Machine dependent opti							
(5) Explain Random Entry Searching. 4M Ans: All Binary Search algorithms, which are fast, but can only operate on tables that are ordered and packed, i.e. tables that will have adjacent items ordered by keywords. Such search procedures may therefore have to be used in conjunction with a sort algorithm which both orders and packs the data. Actually, it is unnecessary for the table to be ordered and packed to achieve good speed in searching. This is also possible to do considerably better with an unpacked, unordered table, provided it is sparse, i.e. the number of storage spaces allocated to it exceeds the number of items to be stored. It is observed that the address calculation sort gives good results with a sparse table. However, having to put elements in order slows down the process. A considerable improvement can be achieved by inserting element in a random (or pscudo-random) way. The random entry number K is generated from the key by methods similar to those used in address calculation. If the K the some other cell must be found for the insertion. The first problem is the generation of a random number from the key. It is to design a procedure that will generate pseudo-random, consistent table positions for keywords. One fairly good prospect for four character EBCDIC keywords is to simply divide the keyword by the table length N and use the remainder. This scheme works well as long as N and the key size (32 bits in case) have no common factors. For a given group of M keywords the remainders should be fairly evenly distributed over) (N-1). 4M (6) Explain Machine dependent optimization phase of compiler with example. 4M Ans: Machine dependent optimization in each explores. 1 • Following figure depicts the matrix that we previ							
Ans: All Binary Search algorithms, which are fast, but can only operate on tables that are ordered and packed, i.e. tables that will have adjacent items ordered by keywords. Such search procedures may therefore have to be used in conjunction with a sort algorithm which both orders and packet to be used in searching. This is also possible to do considerably better with an unpacked, unordered table, provided it is sparse, i.e. the number of storage spaces allocated to it exceeds the number of items to be stored. It is observed that the address calculation sort gives good results with a sparse table. However, having to put elements in order slows down the process. A considerable improvement can be achieved by inserting element in a random (or pseudo-random) way. The random entry number K is generated from the key by methods similar to those used in address calculation. If the K the some other cell must be found for the insertion. The first problem is the generate pseudo-random, consistent table positions for keywords. One fairly good prospect for four character EBCDIC keywords is to simply divide the keyword by the table length N and use the remainder. This scheme works well as long as N and the key size (32 bits in case) have no common factors. For a given group of M keywords the remainders should be fairly evenly distributed over) (N-1). 4M Ans: Machine dependent optimization: 6 • Following figure depicts the matrix that we previously optimized by eliminating a common sub expression (M4). • Next to each matrix erry is a code generated using the operators. • The third column is even better code in that it uses less storage and is faster due to a more appropriate mix of instructions. • The third column is even better code in that it uses less storage and is faster due to a more appropriate m	(5)	Explain Rando	m Entry Searchi	ng.			4M
 Initial of the problem is the series of the end end of the end o	Ans	All Binary Seat	ch algorithms w	hich are fast but	can only operate	e on tables that ar	e (Explanation: 4
(6)Explain Machine dependent optimization phase of compiler with example.4MAns:Machine dependent optimization: • If we optimize register usage in the matrix, it becomes machine – dependent optimization.(Description:2 marks, 		ordered and pac search procedur which both ord ordered and pac considerably be number of stora observed that t However, havin improvement ca way. The rand those used in ac insertion. The fi design a procee keywords. One divide the keyw well as long as given group of 1 (N-1).	1 marks) 1 20 20 21 22 23 24 25 26 27 28 29 20 21 22 23 24 25 26 27 28 29 20 21 22 23 24 25 26 27 28 29 20 21 22 <				
 Ans: Machine dependent optimization: If we optimize register usage in the matrix, it becomes machine – dependent optimization. Following figure depicts the matrix that we previously optimized by eliminating a common sub expression (M4). Next to each matrix entry is a code generated using the operators. The third column is even better code in that it uses less storage and is faster due to a more appropriate mix of instructions. This example of machine-dependent optimization has reduced both the memory space needed for the program and the execution time of the object program by a factor of 2. Machine dependent optimization is typically done while generating code 	(6)	Explain Machi	4M				
	Ans:	 Machine depend If we optimization. Following fig common sub e Next to each the third columnore appropriate this example needed for the the Machine dependent. 	lent optimization: ize register usage gure depicts the n expression (M4). matrix entry is a communication mn is even better ate mix of instruct of machine-depen program and the ndent optimization	e in the matrix, natrix that we pro ode generated usi code in that it us tions. dent optimization execution time of n is typically done	it becomes man eviously optimize ng the operators. es less storage ar has reduced both the object progra	chine – dependented by eliminating and is faster due to a the memory space of the memory space of 2.	t (Description:2 marks, Example: 2 a marks)



MODEL ANSWER

SUMMER- 17 EXAMINATION

Subject Title: System Programming

Subject Code:

		Г		C	ntimized M	latrix		First try		Improved o	ode				
		ŀ		-	punned in	auna	L	1 START	T	1 START					
		ŀ	1		STADT	EINICH	C C	1 EINISU	S	1 EINISU	MI		D 1		
		ŀ	1	-	START	FINISH	O	1, 111131	3	1, FINISH	IVII	-	KI		
		ŀ	-+				51	1, 1911							
		⊦						1.0.475	I.	0.0475	<u> </u>				
		-	-				L	I, RATE	L	3, RATE					
		ļ	2	*	RATE	M1	M	0, M1	MR	2, 1	M2	\rightarrow	R3		
		Ļ					ST	1, M2							
		Ļ													
							L	1, =F'2'	L	5, = F'2'					
			3	*	2	RATE	Μ	0, RATE	M	4, RATE	M3	\rightarrow	R5		
		L					ST	1, M3							
		L													
			4												
		Γ													
		Γ													
		ſ					L	1, M1							
		F	5	-	M1	100	S	1, =F'100'	S	1, =F'100'	M5	\rightarrow	R1		
		ŀ					ST	1, M5							
		ŀ						,					\vdash		
		ŀ					L	1 M3	LR	7.5					
		ŀ	6	*	M3	M5	M	0 M5	MR	6.1	M6	\rightarrow	R7		
		ŀ	·		1415	1415	ST	1 M6	- Mile	0,1		-	K /		
		ŀ	-+				51	1, 140							
		ŀ					T	1 M2							
		ŀ	7	+	M2	M6		1, 112	AD	2 7	M7		D2		
		ŀ	'	т	IVIZ	NIO	A	1, 100	AK	3, 7	1017		K5		
		ŀ	-+				т	1, 1/1	CT.	2 COST					
		ŀ	-+				L	1, 1/1	51	3,0031					
		ŀ	0		M7	COST	eT	1 COST	l						
		L	0		INI /	COST	51	1,0051							
5.		Solv	ve a	ny	FOUR:										16Marks
				·											
	(1)	Exp	lai	n m	nacro witl	nin macro).								4M
		_													
	Ans:	Mac	ro	call	ls are "abl	oreviation	s" of	instruction s	sequen	ices, it seems 1	reasor	nable	e that	such	(Explanation:
		"abb	orev	viat	ions" shoi	uld be ava	ilable	e within othe	r macı	o definitions.					4 marks)
															,
		For	exa	mp	ole,										
				1											
				MA	ACRO										
				ΔГ	D1 & AR	G									
				τι 1		0									
				LI	, &AKG										
				A 1	l,=F"1"										
				ST	1. &ARC	ŕ									
				М		-									
				11/1	DIN L										
				_											
				MA	ACRO AE	DDS &AR	G1, 6	&ARG2, &A	RG3						
				AΓ	DD1 &AR	G1									
				ΔГ	 <i>۹،</i> ۸۵ (G2									
				AL		02									



SUMMER- 17 EXAMINATION

Subject Title: System Programming

Subject Code:





SUMMER-17 EXAMINATION

17634 Subject Code: Subject Title: System Programming Working of absolute loader: With an absolute loading scheme, the programmer and assembler perform the task of allocation, relocation and linking. Therefore, it is only necessary for the loader to read cards of the object deck and move the texts on the cards into the absolute locations specified by the assembler. There are two types of information that the object deck must communicate from the assembler to the loader. First, it must convey the machine instructions that the assembler has created along with the assigned core locations, (called as text cards). Second, it must convey the entry point of the program, which is where the loader is to transfer control when all instructions are loaded (called as transfer cards). Algorithm: • The object deck for this loader consists of a series of text cards terminated by the transfer cards. Therefore, the loader should read one card at a time, moving the text to the • location specified on the card, until the transfer card is reached. At this point the assembled instructions are in core, and it is only necessary to transfer to the entry point specified on the transfer card. (3) Explain general model of a compiler. 4MIn analyzing the compilation of simple program there are seven distinct logical (Figure: 2 Ans: problems as follows and summarized in figure below. marks ,Working: 2 Term marks) Reductions lanent data base Interpret Machine 4) 2) Syntax 1) Lexical tation Storage 6) Code Assembly independent analysis analysis (action assignme selectio and output optimization Uniform Source Relocatable Optimized Assembly Matrix symbol code matrix bject code code table Created data bases Identifie tabi Literal Solid lines indicate creation of data, dashed lines indicate references. table The phases of the compiler communicate by data bases



MODEL ANSWER

SUMMER-17 EXAMINATION

<u>Subj</u> ect	Title: System ProgrammingSubject Code:1763	4
	 Lexical analysis – Recognition of basic elements of creation of uniform symbols. Syntax analysis – Recognition of basic syntactic constructs through reductions. Interpretation – Definition of exact meaning, creation of matrix and tables by action routines. Machine Independent Optimization – Creation of more optimal matrix. Storage Assignment – Modification of identifier and literal tables. If makes entries in the matrix that allow code generation to create code that allocates dynamic storage and that also allow the assembly phase to reserve the proper amounts of STATIC storage. Code Generation – Use of macro processor to produce more optimal assembly code. Assembly And Output – Resolving symbolic addresses and generating machine language. 	
(4)	Mention four functions of storage assignment phase of compiler.	4M
Ans:	 The purpose of this phase is to: Assign storage to all variables referenced in the source program. Assign storage to all temporary locations that are necessary for intermediate result, e.g the results of matrix lines. These storage references were reserved by the interpretation phase and did not appear in the source code. Assign storage to literals. Ensure that the storage is allocated and appropriate locations are initialized (Literals and any variables with the initial attribute) The storage allocation phase first scans through the identifier table, assigning locations to the storage allocation phase first scans through the identifier table, assigning location counter, initialized at zero, to keep track of how much storage it has assigned. Whenever it finds a static variable in the scan, the storage allocation phase does the following four steps: Updates the location counter with any necessary boundary alignment. Calculate the length of the storage needed by the variable (by examining its attributes). 	(Explanation: 4 marks)
	5. This allows code generation to generate the proper amount of storage. For each	



SUMMER-17 EXAMINATION

Subject Title: System Programming

Subject Code: 17634

entry:					
	Initialize	Variable			
6. This tells code values that the action identifier table is mad- relative location for ea made in the matrix. Or address part of instruc- controlled. However, allocates this storage allocates automatic sto	e generation to n routines save le for automation ach entry. An "a Code generation ctions. No stora the matrix entry at execution tim- prage.	put into the d in the iden e storage and automatic" en n use the rel ge is generat y automatic c me, i.e., when	e proper s ntifier tab controlled atry and a ative loca ded at com loes cause n the gene	torage loca le. A simil d storage. T "controlled tion entry pile time for code to be erated code	ation the initial lar scan of the The scan enters "entry are also to generate the or automatic or generated that is executed, it
	Ι	LIT Size			

static area and initializes the storage with the values of the literals. Temporary storage is handled differently since each source statement may reuse the temporary storage (intermediate matrix result area) of the previous source statement. A computation is made of the temporary storage that is required for each source statement. The statement required the greatest amount of temporary storage determines the amount that will be required for the entire program. A matrix entry is made of the form this enables the code generation phase to generate code to create the proper amount of storage.

AUTOMATIC	Size	
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8. Temporary storage is automatic since it is only referenced by the source program and only needed while the source program is active.

 (5)	Explain single pass algorithm for macro processing.	4M
Ans:	If we wanted to provide for macro definitions within macros. The basic problem here is	(Explanation:
	that inner macro is defined only after the outer one has been called in order to provide	4 marks)
	for any use of the inner macro we would have to repeat both the macro-definition and	,
	the macro-call passes. However there is a simpler solution that has added advantages of	
	reducing all macro processing to a single pass. There are two additional variables	
	introduced in the one -pass design a macro definition input (MDI) indicator and a	



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Macro Definition Level Counter (MDLC). The MDI and MDLC are switches (counters) used to keep track of macro calls and macro definitions. The MDI indicator has the value "ON" during expansion of a macro call and the value "OFF" at all other times. The actual expansion of macro calls is performed in the read box. READ tests the switch MDI. If it is "ON", lines are read from the Macro Definition Table (MDT). The reading of a MEND line indicates the end of a macro and terminates expansion of call: MDI is reset to "OFF" and the next line is obtained from the regular input stream. Note that lines returned by READ may include macro definition's; expanded macro code comes out of READ looking just like any other code and may therefore include macro definitions. The macro definition level counter is incremented by 1 when a MACRO pseudo-op is encountered and decremented by 1 when a MACRO pseudo-op occurs. The MDLC is used to insure that the entire macro definition, including MACROs and MENDs, gets stored in MDT. Working of Macro: 1. Start **2.** Initialize MDTC = 1MNTC =1 MDI = OFFMDLC = 03. Read- perform read operation of macro. 4. Search MNT for match found of operation code. 5. Is macro name found ? macro call If yes MDI ='ON' MDTP = MDT Index from MNT entry Setup ALA Goto step 3 If no go to step 6 6. Is macro pseudo-op ?macro definition, If yes MDLC = MDLC + 1If no then read code again. 7. Store macro name and current value of MDTC in MNT entry number in MNTC. **8.** Increment MNTC by 1 Prepare ALA 9. Store macro name card in MDT Increment MDTC by 1 Increment MDLC by 1 10. Read - perform read operation of macro. 11. Substitute index notation for arguments in definition. Enter line is MDT Increment MDTC by 1 **12.** Is MACRO pseudo-op If yes increment MDLC by 1 and goto step 10. If no Is MEND pseudo-op? If yes



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	Decrement MDLC by 1 Goto step 13. If no, goto step 10 13. If MDLC = 0 If yes, goto step 3 If no, goto step 10 14. Write expanded code in source file card. 15. Is END pseudo-op ? If yes Supply expanded source file to assembler processing If no, goto step 3	
(6)	Define parser. Draw the parse tree for string 'abccd' using top-down parser.	4M
Ans:	 Parser: A parser is a program that receives input in the form of sequential source program instructions, interactive online commands, markup tags, ets and breaks them up into parts such as mnemonics, symbols, objects, methods, etc that can then be managed by other phases of compiler. Parser is also called as "Syntax analyzer". Parse tree for the string 'abccd' using top down parser. String is "abccd" Assume: S → xyz aBC B → b bc C → dc cd Steps Assertion 1: abccd matches S Assertion 2: abccd matches approximately and the set of a set of a	(Definition: 1 mark, Parse Tree: 3 marks)
	• Assertion 4 is true.	

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<u>s</u> -	ubject	Title: System Programming Subject Code:	17634	
		 Assertion 3 is true. Assertion 2 is true. Assertion 1 is true. 		
		SSSS $ $ $/ $ $/ $ $/ $ $/ $ a B C b b c c d b b c d b b c d b c c d c c d d c d d d c d <td< th=""><th></th><th></th></td<>		
6.		Solve any FOUR:		16Marks
	(1)	Explain assembly phase of a compiler.		4M
	Ans:	 After code generation phase, next phase is assembly phase. The ta assembly phase depends on how much has been done in code generation. If a lot of work has been done in code generation, then the assembly phase resolve labels references in object program, format the object deck, and the appropriate information for the loader. If code generation has simply generated symbolic machine instruction labels, the assembly phase must (1) resolve label references, (2) ca addresses, (3) generate binary machine instructions, and (4) generate st convert literals. Databases: Identifier table: assembly phase uses this database to enter the value of all into identifier table. Literal table: places the literal on appropriate TXT cards Object code: the output of code generation. Algorithm: A simple assembly phase scans the object code, resolving all label refeared and producing the TXT cards. It hen scans the identifier table to create the ESD cards. 	ask of (se must format ns and lculate torage, l labels erences and the	(2 marks Explanation: 2 marks Databases: , 1 mark, Algorithm: 1 mark)



MODEL ANSWER

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(2)	Sort following numbers in ascending order by bucket sort: 78,354,51,278,63,89,312,12	4M
Ans:		(4 marks)
	0-100 101-200 201-300 78, 51, 63, 278 354 312	
	89, 12 Sort individual buckets	
	12, 51, 63, 78, 89 278 312, 354	
	Sorted list: 12, 51, 63, 78,89, 278, 3	
(3)	Explain features of macro facility.	4M
Ans:	 Following are features of a MACRO facility: 1. Macro instruction Arguments. 2. Conditional Macro Expansion 3. Macro calls within Macros 4. Macro instruction Defining Macros 1. Macro Instruction Arguments: The macro facility presented is capable of inserting block of instructions in place of macro calls. All of the calls to any given macro will be replace by identical blocks. This lacks flexibility: there is no way for a specific macro call to modify the coding that replaces it. An important extension of this facility consists of providing for arguments. An important extension of this facility consists of providing for arguments or parameters in macro calls. Corresponding macro dummy arguments will appear in macro definitions. 	(Each feature: 1 mark)
	· · · A 1,DATA1 A 2,DATA1 A 3,DATA1 · · · A 1,DATA2 A 2,DATA2 A 3,DATA2	



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	· DATA F ^{**} 5 [*] 1 DC DATA F ^{**} 10 2 DC MACRO INCR A A A MEND · · INCR	,)" DATA1	Macro INC Argument &ARG 1,&ARG 2,&ARG 3,&ARG Use ope	R has One e DATA1 as rand		
	· INCR · · DATA1 DC DATA2 DC	DATA2	Use ope F"5" F"10"	e DATA2 as rand		

In this case the instruction sequences are very similar but not identical. The first sequences performs an operation using DATA1 as operand; the second using DATA2. They can be considered to perform the same operation with a variable parameter, or argument. Such parameter is called a macro instruction argument or dummy arguments. It is specified on the macro name line and distinguished by the ampersand which is always its first character.

2. Conditional Macro Expansion: Two important macro processor pseudo-ops, AIF and AGO, permit conditional reordering of the sequence of macro expansion. This allows conditional selection of the machine instructions that appear in expansions of macro call. AIF is conditional branch pseudo-o; it performs an arithmetic test and



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	 branches only if the tested condition is true. The AGO is an unconditional branch pseudo ops or 'go to' statement. It specifies a label appearing on some other statement in the macro instruction definition; the macro processor continues sequential processing of instruction with the indicated statement. These statements are directives to the macro processor and do not appear in macro expansion. 3. Macro call within Macros: Since macro calls are "abbreviations" of instruction sequences, it seems reasonable that such abbreviations should be available within other macro definitions. Macro calls within macros can involve several levels. In fact, conditional macro facilities make it possible for a macro to call itself. So long as this does not cause an infinite loop so long as at some point the macro having been called for the nth time, decides not to call itself again perfectly. 4. Macro Instructions Defining Macros Macros are generalized abbreviations for instruction sequences, nothing that it seems reasonable to permit any valid statements in the abbreviated sequence, including macro definitions. In this manner a single macro instruction might be used to simplify the process of defining a group of similar macros. 	
(4	4) List the steps for binary search algorithm. List the best, worst and average case complexity.	4M
A	Ans: Binary Search Algorithm: A more systematic way of searching an ordered table. This	(Algorithm: 2
	 Find the middle entry (N/2 or (N+1)/2) Start at the middle of the table and compare the middle entry with the keyword to be searched. The keyword may be equal to, greater than or smaller than the item checked. The next action taken for each of these outcomes is as follows If equal, the symbol is found If greater, use the top half of the given table as a new table search If smaller, use the bottom half of the table. Example: The given nos are: 1,3,7,11,15 To search number 11 Indexing the numbers from list [0] upto list[5] 	marks Complexity: 2 marks)
	 technique uses following steps for searching a keywords from the table. 1. Find the middle entry (N/2 or (N+1)/2) 2. Start at the middle of the table and compare the middle entry with the keyword to be searched. 3. The keyword may be equal to, greater than or smaller than the item checked. 4. The next action taken for each of these outcomes is as follows If equal, the symbol is found If greater, use the top half of the given table as a new table search If smaller, use the bottom half of the table. Example: The given nos are: 1,3,7,11,15 To search number 11 Indexing the numbers from list [0] upto list[5] Pass 1 Low=0, High = 5, Mid= (0+5)/2 = 2 	marks Complexity: 2 marks)



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	list[2] = 3 is less than 7	
	$\operatorname{Berr} 2$	
	Pass 2	
	Low= $(Mid+1)/2$ i.e $(2+1)/2 = 1$, High = 5, Mid= $(1+5)/2 = 6/2 = 3$	
	So list [3] = 11	
	and the number if found.	
	List the best, worst and average case complexity:	
	 Binary Search can be analyzed with the best, worst, and average case number of comparisons. These analyses are dependent upon the length of the array, so let N = A denote the length of the Array A. The numbers of comparisons for the recursive and iterative versions of Binary Search are the same, if comparison counting is relaxed slightly. For Recursive Binary Search, count each pass through the if-then-else block as one comparison. For Iterative Binary Search, count each pass through the while block as one comparison. 	
	 Best case - O (1) comparisons In the best case, the item X is the middle in the array A. A constant number of comparisons (actually just 1) are required. Worst case - O (log n) comparisons In the worst case, the item X does not exist in the array A at all. Through each recursion or iteration of Binary Search, the size of the admissible range is halved. This halving can be done ceiling (lg n) times. Thus, ceiling (lg n) comparisons are required. Average case - O (log n) comparisons To find the average case, take the sum over all elements of the product of number of comparisons required to find each element and the probability of searching for that element. To simplify the analysis, assume that no item which is not in A will be searched for, and that the probabilities of searching for each element are uniform. 	
(5)	Describe the term token with respect to Lexical Analysis.	4M
Ans:	 The first phase of compiler is lexical analysis. It works as a text scanner. This phase scans the source code as a stream of characters and converts it into meaningful lexemes. Lexical analyser represents these lexemes in the form of tokens as: <token-name, attribute-value=""></token-name,> Algorithm of Lexical Analysis phase of compiler is as follows The first tasks of the lexical analysis algorithm are to the input character string into token. The second is to make the appropriate entries in the tables. 	(Explanation: 4 marks)
	 A token is a substring of the input string that represents a basic element of the language. It may contain only simple characters and may not include another token. To the rest of the compiler, the token is the smallest unit of currency. Only lexical analysis and the output processor of the assembly phase concern themselves with such elements as characters. Uniform symbols are the terminal symbols for syntax analysis. 	



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	• Lexical analysis recognizes three types of token: Terminal symbols, possible identifiers, and literals.	
	• It checks all tokens by first comparing them with the entries in the terminal table. Once a match is found, the token is classified as a terminal symbol and lexical analysis creates a uniform symbol of type "TRM., and inserts it in the uniform symbol table. If a token is not a terminal symbol, lexical analysis proceeds to classify it as a possible identifier or literal. Those tokens that satisfy the lexical rules for forming identifiers are classified as "possible identifiers".	
	• Example: Consider following program WCM: PROCEDURE(RATE,START,FINISH); DECLARE (COST,RATE,START,FINISH) FIXED BINARY (31)STATIC; COST = RATE * (START-FINISH) + 2*RATE*(START-FINISH-100); RETURN (COST); END:	
	WCM : PROCEDURE (RATE , START , FINISH) ; DECLARE (COST , RATE , START , FINISH) FIXED BINARY (31) STATIC ;	
	COST = RATE * (START - FINISH) + 2 * RATE * (START - FINISH - 100) ; RETURN (COST) ; END ;	
(6)	Explain Conditional Macro expansion.	4M
Ans:	• Two important macro-processor pseudo-ops AIF and AGO permit conditional reordering of the sequence of macro expansion. This allows conditional selection of the machine instructions that appear in expansions of Macro call. Consider the following program.	(Explanation: 4 marks)
	Loop 1 A1, DATA 1 A2, DATA 2 A3, DATA 3	
	Loop 2 A1, DATA 3	
	A2, DATA 2	
	Loop 3 A1, DATA1	



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DATA 1 DC F'5' DATA 2 DC F'10' DATA 3 DC F'15'			
• In the below example, the operands, labels and the nur generated change in each sequence. The program can writte	nber of instructio n as follows:-	ns	
· · · · · · · · · · · · · · · · · · ·	&ARG3		
 LOOP1 VARY 3,DATA1,DATA2,DATA3 LO . A 2,DATA2 . A 3,DATA3	DOP1 A 1,DATA	1	
LOOP2 VARY 2,DATA3,DATA2 LOOP2 A . A 2,DATA2 	1,DATA3		
LOOP3 VARY 1,DATA1 LOOP3 A 1,DATA	1		
DATA1 DC F'5' DATA2 DC F'10' DATA3 DC F'15'			
 Labels starting with a period (.) such as .FINI are appear in the output of the macro processor. The statement AIF (& COUNT EQ1) .FINI direct the the statement. Labelled .FINI if the parameter correspon otherwise the macro processor is to continue with the stat pseudo-ops. AIF is conditional branch pseudo ops it perforbranches only if the tested condition is true. AGO is an unconditional branch pseudo-ops or 'Go label appearing on some other statement. AIF & AGO which the macro processor expands the statements in macro. 	macro labels and macro processor to ding to & COUN tement following trms an arithmetic to' statement. It controls the seq o instructions.	d do not to skip to NT is a1; the AIF test and specifies uence in	